Expected utility and Mixture space theorem

Rohit Lamba Pennsylvania State University rlamba@psu.edu

January 2020

1 Preliminaries and linear representation

Let X be a finite set of prizes. Let $\mathcal{L}(X)$ be the set of lotteries on X, i.e.,

$$\mathcal{L}(X) = \{p : X \to [0,1] | \sum_{x \in X} p(x) = 1\}$$

To simplify notation¹ we will use \gtrsim for a preference relation and > for a *K*-preference relation and also use there for the corresponding derived binary relationships (e.g. $>_{\gtrsim} \equiv >$ etc.). The context should be clear in each case.

For any $x \in X$, δ_x denotes the lottery that chooses x with certainty. That is,

$$\delta_x(y) = \begin{cases} 1 & \text{if } y = x \\ 0 & \text{if } y \neq x \end{cases}$$

Let any $p, q \in \mathcal{L}(X)$ and $a \in [0, 1]$. Then $ap + (1 - a)q \in \mathcal{L}(X)$ and is often referred to as a compound lottery. Dixit writes

A general presumption of standard microeconomic theory is that people care only about what they finally get to consume, and not by the path or process by which they arrived at this consumption vector. This is sometimes called the compound lottery axiom or the reduction of compound lotteries. This axiom will be violated if the decision-maker likes or dislikes the process by which uncertainty is resolved; for example he may enjoy the suspense that builds up as a lottery yields a prize that is another lottery, or he may find it worrying.

In this chapter, we will assume that the compound lottery axiom.

We say that $U : \mathcal{L}(X) \to \mathbb{R}$ is linear if

 $U(ap + (1 - a)q) = aU(p) + (1 - a)U(q) \quad \forall p, q \in \mathcal{L}(X) \text{ and } a \in [0, 1]$

We say that U is a linear representation of \gtrsim if U is linear and represents \gtrsim .

¹Prof. Gul says "The gloves are now off!"

Proposition 1. If U is a linear function on $\mathcal{L}(X)$, then there exists a real-valued function u on X such that $U(p) = \sum_{x \in X} p(x)u(x)$ for all $p \in \mathcal{L}(X)$.

Proof. For any $x \in X$ define $u(x) = U(\delta_x)$. Further, for any $p \in \mathcal{L}(X)$, define $S_p = \{x \in X | p(x) > 0\}$. S_p is basically the support of p. We prove the result by induction on $|S_p|$.

If $|S_p| = 1$, then the claim is trivial. We have $u(x) = U(\delta_x)$ where $\{x\} = S_p$. Suppose we have shown the result for $|S_p| = k$. We claim that it holds for $|S_p| = k + 1$.

Fix any $x \in X$. Note that, any $p \in \mathcal{L}(X)$ can be written as,

$$p = p(x)\delta_x + (1 - p(x))q$$

where

$$q(z) = \begin{cases} 0 & \text{if } z = x \\ \frac{p(z)}{1 - p(x)} & \text{if } z \neq x. \end{cases}$$

Clearly, $q \in \mathcal{L}(X)$ and $|S_q| = k$. Therefore, by the induction hypothesis, $U(q) = \sum_{z \in x} q(z)U(\delta_z)$. Thus,

$$U(p) = p(x)U(\delta_x) + (1 - p(x))U(q) = \sum_{z \in X} p(z)U(\delta_z)$$

and we are done.

If X was an arbitrary srt, the above construction works if we consider $\mathcal{L}(X)$ to be the set of all simple lotteries on X.² Note that even if X is countably infinite and we consider non-simple lotteries, the above construction does not work.

2 von Neumann-Morgenstern Expected Utility

von Neumann-Morgenstern were one of the first to provide a full theory of expected utility in an axiomatic fashion. It is important to note that the vNM model views uncertainty as objective, in the sense that there is given a quantification of how likely the various outcomes are, given in the form of a probability distribution. The other view by Savage regards uncertainty as subjective, which we will discuss later.

Axiom 1: \gtrsim is a preference relation.

Axiom 2: For all $p, q, r \in \mathcal{L}(X)$, p > q and $a \in (0, 1) \implies ap + (1 - a)r > aq + (1 - a)r$.

Axiom 3: For all $p, q, r \in \mathcal{L}(X)$, $p > q > r \implies \exists a, b \in (0, 1)$ such that ap + (1 - a)r > q > bp + (1 - b)r.

²A simple lottery is a lottery with a finite support.

Axiom 2 is often called the independence or substitution axiom, and Axiom 3 is often called the Archimedean or continuity axiom.

Theorem 1. Let X be a finite set and \geq be a binary relation on $\mathcal{L}(X)$. Then, \geq satisfies Axioms 1-3 iff it has a linear representation.

Proof. We break the proof into the following lemmas.

Lemma 1. *If* \geq *satisfies Axioms 1-3, then:*

- (i) p > q and $0 \le a < b \le 1 \implies bp + (1-b)q > ap + (1-a)q$
- (ii) $p \gtrsim q \gtrsim r$ and $p > r \implies \exists a \text{ unique } a^* \in [0, 1] \text{ such that } q \sim a^*p + (1 a^*)r$
- (iii) $p \sim q \text{ and } a \in [0,1] \implies ap + (1-a)r \sim aq + (1-a)r \ \forall r \in \mathcal{L}(X).$

Proof of Lemma 1. (i) First consider the special case a = 0. Then, p > q and $0 \le b \le 1$ with Axiom 2 imply bp + (1-b)q > bq + (1-b)q = q = ap + (1-a)q. Now let r = bp + (1-b)q and suppose a > 0. Then, (a.b) < 1, and r > q and Axiom 2 together imply that

$$r = (1 - \frac{a}{b})r + \frac{a}{b}r$$

> $(1 - \frac{a}{b})q + \frac{a}{b}r$
= $(1 - \frac{a}{b})q + \frac{a}{b}(bp + (1 - b)q)$
= $ap + (1 - a)q$.

(*ii*) Since p > r, part (*i*) ensures that if a^* exists, it is unique. If $p \sim q$, then $a^* = 1$ works. If $q \sim r$, then $a^* = 0$ works. So we need to only consider the case p > q > r. Define

$$a^* = sup\{a \in [0, 1] | q \ge ap + (1 - a)r\}$$

Since a = 0 is in the set, we know we aren't taking supremum over an empty set.

By definition of a^* , if $1 \ge a \ge a^*$, then ap + (1 - a)r > q. Moreover by (*i*), if $0 \le a \le a^*$, then q > ap + (1 - a)r. To see this, note that if $0 \le a \le a^*$, then there exists *a*' such that $0 \le a \le a' \le a^*$ and $q \ge a'p + (1 - a')r$ by definition of a^* . And a < a' implies that $q \ge a'p + (1 - a')r > ap + (1 - a)r$. There are three possibilities to consider

There are three possibilities to consider.

Suppose $a^*p + (1-a^*)r > q > r$. Then by Axiom 3, there exists $b \in (0, 1)$ such that $b(a^*p + (1-a^*)r) + (1-b)r > q$, that is, $ba^*p + (1-ba^*)r > q$. But $ba^* < a^*$, so we have $q > ba^*p + (1-ba^*)r$. Contradiction.

Suppose $p > q > a^*p + (1 - s^*)r$. Then, by Axiom 3, there exists $b \in (0, 1)$, such that $q > b(a^*p + (1 - a^*)r) + (1 - b)p$, that is, $q > (1 - b(1 - a^*))p + b(1 - a^*)r$. Since, $(1 - b(1 - a^*)) > a^*$, we have from the above argument that $(1 - b(1 - a^*))p + b(1 - a^*)r > q$. Contradiction.

This leaves us with the third possibility (which is what we want), namely $q \sim a^*p + (1 - a^*)r$. *(iii)* This result is trivial for the case where, for all $s \in \mathcal{L}(X)$, $p \sim q \sim s$. So suppose there is some $s \in \mathcal{L}(X)$ such that $s > p \sim q$. (The other case where there is some s such that $p \sim q > s$ can be similarly done.) Let any $r \in \mathcal{L}(X)$ and $a \in [0,1]$. We want to show that $ap + (1 - a)r \sim aq + (1 - a)r$. Suppose ap + (1 - a)r > aq + (1 - a)r. From Axiom 2 we can deduce that for all $b \in (0,1)$, $bs + (1 - b)q > bq + (1 - b)q = q \sim p$. Using Axiom 2 again, we get that a(bs + (1 - b)q) + (1 - a)r > ap + (1 - a)r, for all $b \in (0,1)$. Since (by assumption) ap + (1 - a)r > aq + (1 - a)r, Axiom 3 implies that for each b there exists some $a^*(b)$ such that $ap + (1 - a)r > a^*(b)(a(bs + (1 - b)q) + (1 - a)r) + (1 - a^*(b))(aq + (1 - a)r)$. Fix, say b = 1/2, and let $a^*(1/2)$ be written a^* ; then we have that

$$ap + (1-a)r > [a^*a/2]s + [a^*a/2 + (1-a^*)a]q + [1-a]r$$

But the term on the right hand side is

$$a[(a^*/2)s + (1 - a^*/2)q] + (1 - a)r$$

and since $a^*/2$, this must be > ap + (1-a)r, a contradiction. The case where we start by assuming aq + (1-a)r > ap + (1-a)r will similarly lead to a contradiction. Thus, we must have $ap + (1-a)r \sim aq + (1-a)r$.

Lemma 2. If \succeq on $\mathcal{L}(X)$ satisfies Axioms 1-3, then there exists \overline{z} and \underline{z} in X such that $\delta_{\overline{z}} \succeq p \succeq \delta_{\underline{Z}}$ for all $p \in \mathcal{L}(X)$.

Proof. Since X is finite, there exists \overline{z} and \underline{z} in X such that $\delta_{\overline{z}} \geq \delta_x \geq \delta_{\underline{z}}$ for all $x \in X$. Let any $p \in \mathcal{A}(X)$. We prove this by induction on the cardinality of S_p , where $S_p = \{x \in X | p(x) > 0\}$.

If $|S_p| = 1$, then $p = \delta_x$ for some $x \in X$, and we already know $\delta_{\overline{z}} \gtrsim \delta_x \gtrsim \delta_{\underline{z}}$. Suppose we have shown the result for $|S_p| = k$. We claim that it also holds for $|S_p| = k + 1$.

Fix any $x \in X$. Note that p can be written as,

$$p = p(x)\delta_x + (1 - p(x))q$$

where

$$q(z) = \begin{cases} 0 & \text{if } z = x \\ \frac{p(z)}{1 - p(x)} & \text{if } z \neq x \end{cases}$$

Clearly, $q \in \mathcal{L}(X)$ and $|S_q| = k$. Therefore, by induction hypothesis, $\delta_{\bar{z}} \gtrsim q$. Now $\delta_{\bar{z}} \gtrsim \delta_x \implies \delta_{\bar{z}} > \delta_x$ or $\delta_{\bar{z}} \sim \delta_x$, and $\delta_{\bar{z}} \gtrsim q \implies \delta_{\bar{z}} > q$ or $\delta_{\bar{z}} \sim q$. Then using Axiom 2 and lemma 1(*iii*) we can get

$$\delta_{\bar{z}} = p(x)\delta_{\bar{z}} + (1 - p(x))q$$
$$\gtrsim p(x)\delta_x + (1 - p(x))q$$
$$= p$$

Thus $\delta_{\bar{z}} \gtrsim p$. Similarly we can show that $p \gtrsim \delta_z$.

Now we shall use the above lemmas to prove the main theorem. Suppose \geq satisfies Axioms 1-3. Apply lemma 2 to produce $\delta_{\bar{z}}$ and $\delta_{\underline{z}}$. If $\delta_{\bar{z}} \sim \delta_{\underline{z}}$, then set $U(p) = \pi$ for all $p \in \mathcal{L}(X)$, and we are done. So for the rest of the proof we consider the case $\delta_{\bar{z}} > \delta_{\underline{z}}$. For $p \in \mathcal{L}(X)$, define,

$$U(p) = a$$
 where $a\delta_{\bar{z}} + (1-a)\delta_{\underline{z}} \sim p$

By lemma 1(ii) such an *a* exists and is unique, so *U* is well defined. We claim that this *U* will do the job for us.

For all $p, q \in \mathcal{L}(X)$ and $a \in [0, 1]$, by applying lemma 1(iii) twice, we get

$$ap + (1 - a)q \sim a[U(p)\delta_{\bar{z}} + (1 - U(p))\delta_{\underline{z}}] + (1 - a)[U(q)\delta_{\bar{z}} + (1 - U(q))\delta_{\underline{z}}]$$
$$= [aU(p) + (1 - a)U(q)]\delta_{\bar{z}} + [1 - aU(p) - (1 - a)U(q)]\delta_{z}$$

Thus by definition of U,

$$U(ap + (1 - a)q) = aU(p) + (1 - a)U(q).$$

This shows that U is linear.

Finally we want to show that U represents \geq . We claim that

$$U(p) \ge U(q) \text{ iff } U(p)\delta_{\bar{z}} + (1 - U(p))\delta_{z} \gtrsim U(q)\delta_{\bar{z}} + (1 - U(q))\delta_{z}$$

If U(p) = U(q), then the two lotteries are equal. Suppose U(p) > U(q). Then by lemma 1(i), we have $U(p)\delta_{\bar{z}} + (1 - U(p))\delta_{z} > U(q)\delta_{\bar{z}} + (1 - U(q))\delta_{z}$.

Conversely, suppose $U(p)\delta_{\bar{z}} + (1 - U(p))\delta_{\underline{z}} \gtrsim U(q)\delta_{\bar{z}} + (1 - U(q))\delta_{\underline{z}}$. If we were to have U(q) > U(p), then by lemma 1(i) would give us $U(q)\delta_{\bar{z}} + (1 - U(q))\delta_{\underline{z}} > U(p)\delta_{\bar{z}} + (1 - U(p))\delta_{\underline{z}}$, which is a contradiction. Thus, we must have $U(p) \ge U(q)$, which completes our claim.

Now suppose \geq has a linear representation. That is, there exists $U: X \to \mathbb{R}$ such that

$$U(ap + (1 - a)q) = aU(p) + (1 - a)U(q) \quad \forall p, q \in \mathcal{L}(X) \text{ and } a \in [0, 1]$$

and $p \gtrsim q \text{ iff } U(p) \ge U(q) \quad \forall p, q \in \mathcal{L}(X)$

We want to show that Axioms 1-3 hold.

In the following discussion, let p, q, r be any elements in $\mathcal{L}(X)$.

Consider p and q. Clearly, $U(p) \ge U(q)$ or $U(q) \ge U(p)$, that is, $p \ge q$ or $q \ge p$. Next, let $p \ge q$ and $p \ge r$. Then, $U(p) \ge U(q) \ge U(r)$. Thus, $U(p) \ge U(r)$ which implies $p \ge r$. Thus \ge is complete and transitive and hence a preference relation. Axiom 1 is satisfied.

Next suppose p > q and let $a \in (0, 1)$. Then U(p) > U(q). Now, U(ap + (1 - a)r) = aU(p) + (1 - a)U(r) and U(aq + (1 - a)r) = aU(q) + (1 - a)U(r). Also, U(p) > U(q) implies

aU(p) + (1-a)U(r) > aU(q) + (1-a)U(r), that is U(ap + (1-a)r) > U(aq + (1-a)r). Hence, ap + (1-a)r > aq + (1-a)r which satisfies Axiom 2.

Finally, suppose p > q > r. Then, U(p) > U(q) > U(r). Note that $t \in [0, 1]$,

$$tU(p) + (1-t)U(r) > U(q) \text{ iff } t > \frac{U(q) - U(r)}{U(p) - U(r)}, \text{ and}$$

$$tU(p) + (1-t)U(r) < U(q) \text{ iff } t < \frac{U(q) - U(r)}{U(p) - U(r)}$$

Therefore, define

$$a = \frac{U(p) - U(r) + \epsilon}{U(p) - U(r)} \text{ and } b = \frac{U(q) - U(r) - \epsilon}{U(q) - U(r)}$$

where ϵ is chosen small enough so that $a, b \in (0, 1)$. This will do the job. Thus, Axiom 3 is also satisfied.

Theorem 2. If $U : \mathcal{L}(X) \to \mathbb{R}$ represents \geq , then a function $V : \mathcal{L}(X) \to \mathbb{R}$ represents \geq iff there exists real numbers c > 0 and b such that

$$V(p) = cU(p) + d \ \forall p \in \mathcal{L}(X)$$

Proof.

3 Mixture Space Theorem

This section is based mostly on the classic paper by Herstein and Milnor [1953].

Definition 1. Let Π be any set with two operations, viz. \oplus and \circ , defined on it such that for eery $p, q \in \Pi$ and $a, b \in [0, 1]$, we have $a \circ p \oplus (1 - a) \circ q \in \Pi$, and the following properties are satisfied:

(1)
$$1 \circ p \oplus (1-1) \circ q = p$$
.

(2)
$$a \circ p \oplus (1-a) \circ q = (1-a) \circ q \oplus a \circ p$$
, and

(3)
$$b \circ [a \circ p \oplus (1-a) \circ q] \oplus (1-b) \circ q = (ab) \circ p \oplus (1-ab) \circ q$$
.

Then, (Π, \oplus, \circ) *is called a Mixture Space.*

Note that \oplus corresponds to the addition operation of real numbers and \circ corresponds to the multiplication operation of the real numbers. In fact the most obvious examples of the mixture space would be any convex subset *C* of \mathbb{R} , that is (*C*, +, .). In what follows we write '+' for \oplus and '.' for \circ . It makes life simple and the context should be obvious.³

³Then the three properties can be written as

⁽¹⁾ 1p + 0q = p,

⁽²⁾ ap + (1 - a)q = (1 - a)q + ap, and

⁽³⁾ b[ap + (1-a)q] + (1-b)q = (ab)p + (1-ab)q.

Remark 1. Note that if for a mixture space the following hold:

$$ap + (1-a)p = p \tag{(*)}$$

and

$$b[ap + (1-a)q] + (1-b)[cp + (1-c)q] = [ab + (1-b)c]p + [b(1-a) + (1-b)(1-c)]q \quad (\bigstar)$$

where of course $p, q \in \Pi$, and $a, b, c \in [0, 1]$.

Condition (*) is straightforward. ap + (1-a)p = a[1p+0p] + (1-a)p = a[0p+1p] + (1-a)p = 0p+1p = 1p+0p = p where the first equality follows from part (1) of the definition, second follows from part (2), third follows from part (3), the fourth follows from part (2) of the definition while the last follows from part (1).

Exercise 1. Show that (\star) holds.

Let $(\Pi, +, .)$ be a mixture space and \geq is binary relation on Π . Consider the following axioms:

Axiom 1' : \gtrsim is a preference relation.

Axiom 2': For all
$$p, q, r \in \Pi$$
, $p > q$ and $a \in (0, 1) \implies ap + (1 - a)r > aq + (1 - a)r$.
Axiom 3': For all $p, q, r \in \Pi$, $p > q > r \implies \exists a, b \in (0, 1)$ such that $ap + (1 - a)r > q > bp + (1 - b)r$.

As in the vNM model to prove the representation theorem we will need the help of some lemmas. So we shoot the following:

Lemma 3. If \gtrsim satisfies Axioms 1' - 3', then:

(i)
$$p > q$$
 and $0 \le a < b \le 1 \implies bp + (1-b)q > ap + (1-a)q$

(ii)
$$p \gtrsim q \gtrsim r$$
 and $p > r \implies \exists a \text{ unique } a^* \in [0, 1] \text{ such that } q \sim a^*p + (1 - a^*)r$

(iii)
$$p \sim q$$
 and $a \in [0, 1] \implies ap + (1 - a)r \sim aq + (1 - a)r \quad \forall r \in \Pi$.

Proof.

Theorem 3. Suppose Π is a mixture space and \geq is a binary relation on Π . then \geq satisfies Axioms 1' - 3' iff it has a linear representation.

Proof.

The main difference in the proof here will be to get around the problem of best and worst prize that used in Lemma 2.

Remark 2. Note that in theorem 1, we found a linear representation when X is finite. Now let X be an arbitrary set and $\mathcal{L}^{0}(X)$ the set of simple lotteries over X. Then, the set $\mathcal{L}^{0}(X)$ with the usual addition and multiplication forms a mixture space. Thus using the Mixture Space theorem, it is easy to see that the vNM expected utility result can be extended to an arbitrary set and the space of simple lotteries. A binary relation \geq satisfying our favourite axiom has a linear representation of the form $U : \mathcal{L}^{0}(X) \to \mathbb{R}$. Moreover, note that in proving proposition 1 we did not use the finiteness of X, only the fact that the support of the lotteries on X is finite. Thus, the same proposition holds for any arbitrary X and $\mathcal{L}^{0}(X)$. Therefore, we have the following result:

Theorem 4. Let X be an arbitrary set and $\mathcal{L}^{0}(X)$ the set of simple lotteries over X. A binary relation \gtrsim satisfies Axioms 1-3 iff there exists a function $U : \mathcal{L}^{0}(X) \to \mathbb{R}$ that represents \gtrsim . Also there exists $u : X \to \mathbb{R}$ such that $U(p) = \sum_{x \in X} p(x)u(x)$ for all $p \in \mathcal{L}^{0}(X)$.